

Nonuniform FFTs at Flatiron

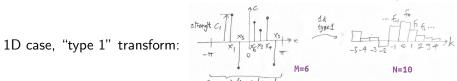
 $-- \sim * OR * \sim -$

Lessons from developing a small numerical library

Alex Barnett¹, with much help from: Jeremy Magland¹, Ludvig af Klinteberg (Mälardalen U.), Melody Shih (NVidia), Joakim Andén (KTH), Libin Lu¹, Robert Blackwell (SCC), Andrea Malleo (Bloomberg), and many others...

FWAM5 Friday, October 20, 2023

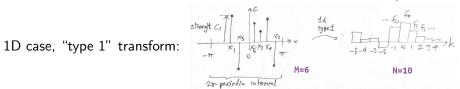
¹Center for Computational Mathematics, Flatiron Institute



Given locations x_1, \ldots, x_M , and their "strengths" c_1, \ldots, c_M , return

$$f_k = \sum_{i=1}^{M} e^{ikx_i} c_i$$
, for $k = -\frac{N}{2}, -\frac{N}{2} + 1, \dots, \frac{N}{2} - 2, \frac{N}{2} - 1$

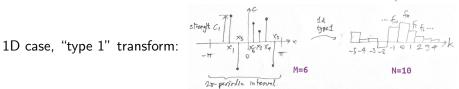
- Has N outputs, each a sum of M terms: naively $\mathcal{O}(NM)$ cost (flops)
- Off-grid version of fast Fourier transform (FFT) there M = N and $x_i = 2\pi j/N$



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- approximates it to a user-requested ε $\varepsilon = 10^{-1}$ (fastest) to $\varepsilon = 10^{-14}$ (slowest)



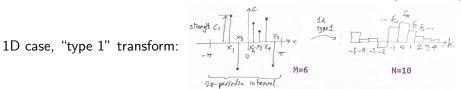
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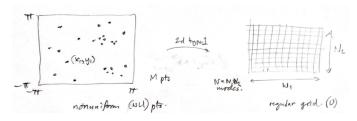
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"Type 3" as type 1 but arbitrary target freqs. $\{s_k\}_{k=1}^N$: $f_k = \sum_{i=1}^M e^{is_k x_j} c_j$ $NU \rightarrow NU$

Higher dimensions also needed

2D type 1:
$$f_{k,\ell} = \sum_{j=1}^{M} e^{i(kx_j + \ell y_j)} c_j$$

for k, ℓ in a rectangle of modes



3D type 1:
$$f_{k,\ell,m} = \sum_{j=1}^{M} e^{i(kx_j + \ell y_j + mz_j)} c_j$$
 etc

- dimensions $\{1,2,3\} \times \text{types } \{1,2,3\} = 9 \text{ transforms}$ Software design: how to avoid code repetition?
- 9 transforms \times {float,double} = 18 functions How reduce number of functions to write and maintain?



Who uses such transforms?

- 1) Fourier image reconstruction: $f_{k,\ell}$ is unknown pixel intensities apparatus measures strengths c_j at nonuniform frequency points (x_j, y_j)
 - MRI (either 2D slice, or 3D)
 - coherent diffraction/powder imaging (X-ray)
 - very long baseline interferometry (VLBI)
 - cryo electron microscopy



some 2D MRI points



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- 2) Numerical forward solvers, simulation
 - electrostatics or fluid problems in periodic box spectral Ewald method: Poisson solve trivial in Fourier space
 - numerical PDE eg, interpolating between overlapping grids
 - eval Fourier transform by numerical quadrature (type 1)
 - Fresnel diffraction (optics)



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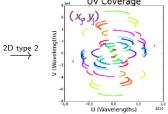
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- 3) Spatial/temporal statistics
 - power-spectrum of NU time-series, or point-masses galaxies
 - fast kernel apply in Gaussian process regression



Probably the most famous astro image of 2019:





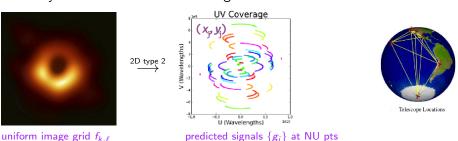




uniform image grid $f_{k,\ell}$

predicted signals $\{g_j\}$ at NU pts

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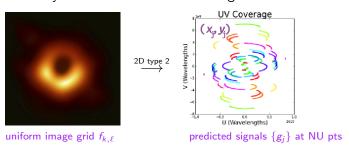


 10^{-10} radian resolution!

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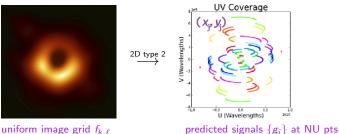


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How? the above is a linear forward model: $\mathbf{g} = A\mathbf{f}$ big dense matrix A lterative optimization of \mathbf{f} until it best fits the detected signal:

$$\mathbf{f}_{\text{recon}} = \operatorname{arg\,min}_{\mathbf{f}} \|A\mathbf{f} - \mathbf{g}_{\text{detected}}\|_{2}^{2} + \lambda_{1} \|\mathbf{f}\|_{1} + \lambda_{\text{TV}} \|\mathbf{f}\|_{\text{TV}}$$

- ullet each iteration, A and A^* applied fast by NUFFTs dominant cost, I think
- same idea in other 2D or 3D Fourier imaging (MRI, cryo-EM, etc)

(Dutt–Rokhlin '93, Steidl '98, Greengard–Lee '04, Potts et al, \dots)

Eg, 1D Type 1:

set up new grid on $[0,2\pi)$ with n=2N points, say

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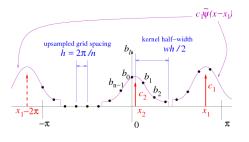
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1) **Spread** each spike c_j onto grid $b_\ell = \sum_{j=1}^M c_j \psi(\ell h - x_j)$ detail: 2π -periodize



 $\mathcal{O}(w^d M)$ flops $\mathcal{O}(wdM)$ kernel evals often dominates cost

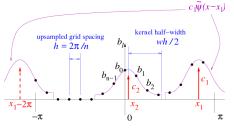
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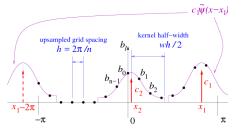
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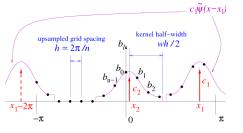
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- 2) Let $\{\hat{b}_k\}$ = size-n **FFT** of $\{b_l\}$
- 3) **Correct** for spreading: $\tilde{f}_k = \frac{1}{\hat{\psi}(k)}\hat{b}_k$ keep only low modes $-N/2 \le k < N/2$
- Type 2 reverses the steps; Type 3 is "Type 2 wrapped inside a Type 1"

fiNUFFT

Prehistory (2015): Leslie Greengard, Jeremy Magland,

Marina Spivak, myself, at SCDA NUFFT needed for imaging (cryo-EM, etc)

had: NYU single-threaded Fortran, Gaussian kernel $\psi(x)$ too wide (Greengard–Lee '09) also: NFFT3 C++, multithreaded, hard to use, user chooses $\psi(x)$ (Keiner et al '06)

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Lesson: keep a CHANGELOG in your repo commit messages not enough ©

"Small" (lib = 3k lines of C++) Dive into some aspects. . .



Performance: kernel $\psi(x)$ evaluation

Problem: for some chips & compilers, $\exp(x)$ slow (40 M evals/sec/core) Fix: piecewise polynomial approx + Horner's rule C gen. by MATLAB!

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```
// Code generated by gen_all_horner_C_code.m in finufft/devel
// Authors: Alex Barnett & Ludvig af Klinteberg.
// (C) The Simons Foundation, Inc.
  if (w==2) {
    FLT c0[] = \{4.5147043243215315E+01, 4.5147043243215300E+01, 0.000000000000E+00, 0.00000000000000E+00\};
    FLT c3[] = \{-2.0382426253182082E+01, 2.0382426253182086E+01, 0.0000000000000E+00, 0.000000000000E+00\};
    FLT C4[1 = \{-2.0940804433577420E+00, -2.0940804433577389E+00, 0.0000000000000E+00, 0.0000000000000E+00\};
    for (int i=0: i<4: i++) ker[i] = c0[i] + z*(c1[i] + z*(c2[i] + z*(c3[i] + z*(c4[i]))):
    FLT c0[] = \{1.5653991189315119E+02, 8.8006872410780295E+02, 1.5653991189967152E+02, 0.0000000000000000E+00\};
   FLT c1[] = {3.1653018869611077E+02, 7.4325702843759617E-14, -3.1653018868907071E+02, 0.00000000000000000E+00};
    FLT c2[1] = \{1.7742692790454484E+02. -3.3149255274727801E+02. 1.7742692791117119E+02. 0.000000000000000E+00\}
    FLT c3[] = \{-1.5357716116473156E+01, 9.5071486252033243E-15, 1.5357716122720193E+01, 0.0000000000000000E+00\};
    FLT C4[] = \{-3.7757583061523668E+01.5.3222970968867315E+01.-3.7757583054647384E+01.0.0000000000000000E+00\};
    FLT C5[] = \{-3.9654011076088804E+00.1.8062124448285358E-13.3.9654011139270540E+00.0.000000000000000E+00\};
    for (int i=0; i<4; i++) ker[i] = c0[i] + z*(c1[i] + z*(c2[i] + z*(c3[i] + z*(c4[i] + z*(c5[i])))));
  } else if (w==4) {
    FLT c0[] = {5.4284366850213200E+02. 1.0073871433088398E+04. 1.0073871433088396E+04. 5.4284366850213223E+02};
    FLT c1[] = {1.4650917259256939E+03. 6.1905285583602863E+03. -6.1905285583602881E+03. -1.4650917259256937E+03};
   FLT c2[] = \{1.4186910680718345E+03, -1.3995339862725591E+03, -1.3995339862725598E+03, 1.4186910680718347E+03\};
    FLT C3[] = \{5.1133995502497419E+02. -1.4191608683682996E+03. 1.4191608683682998E+03. -5.1133995502497424E+02]
    FLT c4[] = {-4.8293622641174039E+01, 3.9393732546135226E+01, 3.9393732546135816E+01, -4.8293622641174061E+01};
    FLT c5[] = {-7.8386867802392288E+01, 1.4918904800408930E+02, -1.4918904800408751E+02, 7.8386867802392359E+01};
    FLT C6[] = \{-1.0039212571700894E+01, 5.0626747735616746E+00, 5.0626747735625512E+00, -1.0039212571700640E+01\};
    for (int i=0; i<4; i++) ker[i] = c0[i] + z*(c1[i] + z*(c2[i] + z*(c3[i] + z*(c4[i] + z*(c5[i] + z*(c6[i])))));
  } else if (w==5) {
    FLT c0[] = {9.9223677575398392E+02, 3.7794697666613320E+04, 9.8715771010760494E+04, 3.7794697666613283E+04, 9.9223P
-:-- ker horner allw loop.c Top L1 Git-master (C/*l WS Abbrev)
```

- GCC/ICC compilers SIMD-vectorize this; get 400-700 M evals/sec/core
- think hard re SIMD, but avoid maintaining intrinsics immintrin.h
- are exploring custom AVX512 (Wenda Zhou + R. Blackwell)



Performance: spreading

The *order* in which NU points spread to grid has big effect on speed!

- bin-sort NU pts (into $16 \times 4 \times 4$ cuboids of grid)
- process all pts in bin 1, then bin 2, ...good for keeping grid in cache(J. Magland)
- multithreaded the bin-sort (2023: M. Reinecke speeds it up!)

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Load-balanced multithreading

Type-2 easy: OpenMP parallel over NU pts no collisons reading from U blocks Type-1 not so: writes collide!

(also collab with J. Magland) load-balance via "subproblems" each of 10⁴ NU pts

 $N_{\rm f}$ 2D case, type-1, spread to fine grid: NU pts x outer prod spread subproblems: each own thread

2017: I wanted simple, familiar to users in C/Fortran and match NYU code

```
C-compatible: pass pointers, explicit array sizes, return value is error code...
int M = 1e8;
                                     // number of nonuniform points
vector<double>
                    x(M):
                                   // NU pts
vector<complex<double> > c(M);
                                    // NU strengths
// (here user fills x and c as they like)
int N = 1e7:
                                     // number of output modes
int status = finufft1d1(M, &x[0], &c[0], +1, 1e-9, N, &f[0], NULL); // do it
```



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• Why last arg NULL? it accepts ptr to options C-struct eg, opts.debug=2



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2018: Batch transforms sharing NU pts: finufft1d1many(ntrans, M, ...)
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batch FFTW often faster, sorting only done once



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2020: maintaining 18 × 2 functions too much pain → "guru" interface 4-function pattern: Create plan, Set the NU pts, Execute transform(s), Destroy plan. finufft_plan: "opaque" pointer to (private) C++ struct. (as in Brian's talk)



2017: I wanted simple, familiar to users in C/Fortran and match NYU code C-compatible: pass pointers, explicit array sizes, return value is error code... int M = 1e8: // number of nonuniform points vector<double> x(M): // NU pts vector<complex<double> > c(M); // NU strengths // (here user fills x and c as they like) int N = 1e7: // number of output modes vector<complex<double> > f(N); // allocate output array int status = finufft1d1(M, &x[0], &c[0], +1, 1e-9, N, &f[0], NULL); // do it Why last arg NULL? it accepts ptr to options C-struct eg. opts.debug=2 2018: Batch transforms sharing NU pts: finufft1d1many(ntrans, M, ...) batch FFTW often faster, sorting only done once 2020: maintaining 18×2 functions too much pain \rightarrow "guru" interface 4-function pattern: Create plan, Set the NU pts, Execute transform(s), Destroy plan. finufft_plan: "opaque" pointer to (private) C++ struct. (as in Brian's talk) Lesson: learn about public vs private headers, namespacing (as I had to)

Lesson: think hard about interface, break it VERY rarely!

help users: preserve all simple and batch interfaces (they call guru)

FLATIRON

N S T I T U T

Wrappers to other languages: expands user base

Guru interface made wrapping easier: just wrap 4 funcs, pass opaque ptr write the simple and batched functions via a few lines in each high-level language



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MATLAB/Octave

(Libin Lu)

- simple for users f = finufft1d1(x,c,+1,1e-9,N);
- still recommend MWrap to auto-gen MEX code horror (Bindel '09)

Python

- 2018: Dan Foreman-Mackey wrote a pybind11 wrapper to FINUFFT
 - served well for a while. But eg forced recompile of libfinufft.so
- 2020: Joakim Andén switched us to ctypes nice, 1000 lines incl. auto-doc-gen
 - Others have wrapped (cu)FINUFFT in autodiff frameworks: tensorflow-nufft, jax-finufft, pytorch-finufft



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Julia wrapper is separate repo: helps separate concerns (L. af Klinteberg)

Lessons: each new language brings installation troubles linux/OSX/Windows



Need tests that check accuracy for all transforms — fail if measured error $\geq 10\varepsilon$?

- > make test [includes bunch of edge cases M=0, N=0, eps=0.0, etc...]
- O segfaults out of 8 tests done
- O fails out of 8 tests done



test 1d type 1:

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A story: I wrote finufft1d_test, etc, writes to stdout for humans...

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1000000 NU pts to 1000000 modes in 0.0829 s 1.21e+07 NU pts/s one mode: rel err in F[370000] is 6.59e-08 [...]

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3 years later: OSX, Windows users cannot find numdiff, ugh! thus: each C++ test driver now uses exit code 0 for success, 1 for fail
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```
Continuous Integration (CI)
```

each GitHub push reruns tests: O C++ Build passing

FLATIRON

Jenkins (was using makefile, now uses CMake)

Documentation

Next topic

This Page

Quick search

README.md, or a .tex file to PDF, fine for small project, but:



FINUFFT is a library to compute efficiently the three most common types of nonuniform fast Fourier transform (NUFFT) to a specified precision, in one, two, or three dimensions, either on a multi-core shared-memory machine, or on a GPU. It is extremely fast (typically achieving 106 to 108 points per second on a CPU, or up to 109 points per second on a GPU), has very simple interfaces to most major. numerical languages (C/C++, Fortran, MATLAB, octave, Python, and Julia), but also has more advanced (vectorized and "guru") interfaces that allow multiple strength vectors and the reuse of FFT plans, The CPU library is written in C++ (with limited use of ++ features), OpenMP, and calls FFTW. It has been developed since 2017 at the Center for Computational Mathematics at the Flatiron Institute by Alex Barnett and others, and is released under an Apache v2 license.

What does FINUFFT do?

As an example, given M real numbers $x_i \in [0, 2\pi)$, and complex numbers c_i , with $i = 1, \dots, M$, and a requested integer number of modes N. FINUFFT can efficiently compute the 1D "type 1" transform, which means to evaluate the N complex outputs

$$f_k = \sum_{i=1}^M c_j e^{ikx_j} \;, \qquad ext{ for } k \in \mathbb{Z}, \quad -N/2 \leq k \leq N/2 - 1 \;.$$

As with other "fast" algorithms. FINUFFT does not evaluate this sum directly-which would take O(NM) effort—but rather uses a sequence of steps (in this case, optimally chosen spreading, EET, and deconvolution) to approximate the vector of answers (1) to within the user's desired relative tolerance, with only $O(N \log N + M)$ effort, ie, in almost linear time. Thus the speed-up is similar to that of the FFT. You may now want to jump to guickstart, or see the definitions of the other transforms in

One interpretation of (1) is: the returned values f_k are the Fourier series coefficients of the distribution $f(x) := \sum_{i=1}^{M} c_i \delta(x-x_i)$, a sum of point-masses with arbitrary locations x_i and strengths c . Such exponential sums are needed in many applications in science and engineering, in-

Web-facing essential for professionalism

- math definitions
- installation
- examples
- tutorial cases
- troubleshooting
- devnotes



 Decent solution I recommend: sphinx, host at readthedocs.io docs/*.rst ReStructuredText, in master branch (unlike GHpages)
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- Always update docs at the time you change the source



GPU version

cuFINUFFT: CUDA called from C++ (Melody Shih, intern '18, '19 [NYU \rightarrow NVidia])

Single-precision V100 GPU comparisons (including H2D & precomp, smaller is better): 2D Type 1, $N_i = 1000$, $M = 10^7$ 3D Type 1, $N_i = 100$, $M = 10^7$ 10^{3} CUNFFT gpuNUFFT ns per nonuniform point CUNFFT gpuNUFFT 10^{2} FINUFFT FINUFFT (28 threads xeon Broadwell) 10^{1} **cuFINUFFT CUFINUEFT** 10^{-4} 10^{-3} 10^{-2} 10^{-1} 10^{-5} 10^{-4} 10^{-3} 10^{-2} 10^{-1} ϵ (relative l_2 error) ϵ (relative l_2 error)

spreading done in shared memory (subproblems), not global memory
 (Shih et al, PDSEC 2021, best paper prize)

2023: cleaned up GPU lib (+tests,...) into FINUFFT repo (R. Blackwell)

Future: language interfaces auto-detect GPU arrays?



First identify a need

+ benchmark existing codes, on same task!

– can you perform task faster / more accurate?



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- users find bugs you'd never thought of must respond (thanks Libin)
- users make improvement PRs, or beat your speed respond to some
- be conservative: one change at a time for your sanity (eg: no MPI)
- try to be somewhat idiomatic | I won't use #define FLT double in C++ again ☺



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Collect a great team: software engineers + interns + motivated users

- library becomes used © ⇒ more Issues / requests ©
- quarterly(?) all-team meeting
 prioritize features, drop others



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Installation, dependencies & language wrappers takes much of our time!

- helps users, but needs a team (linux,OSX,Windows). Your choice.

Eg: guess which language wrapper SMILI black hole astro code uses?

- still plenty of fun with math & floating-point speed hacking!

Thank-you!

